

Efficient GPU Offloading with OpenMP for a Hyperbolic Finite Volume Solver on Dynamically Adaptive Meshes

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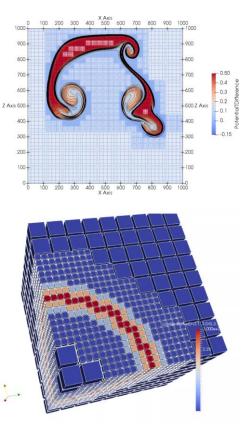


Motivation and Project Overview

• ExaHyPE: A wave equation solver with explicit time stepping

$$\frac{\partial Q}{\partial t} + \nabla \cdot F(Q) + \sum_{i=1}^{d} B_i(Q) \frac{\partial Q}{\partial x_i} = S(Q)$$

- Peano: Dynamically adaptive mesh refinement
- Patches of cells constructed through octree-type AMR
- MPI+X paradigm: Many ranks, each hosting several threads
- Each thread offloads Finite Volume patches to the GPU
- Gather multiple patches into batch of update tasks
- Handle tasks in one rush on the GPU via one kernel call
- Realisation by OpenMP's target constructs

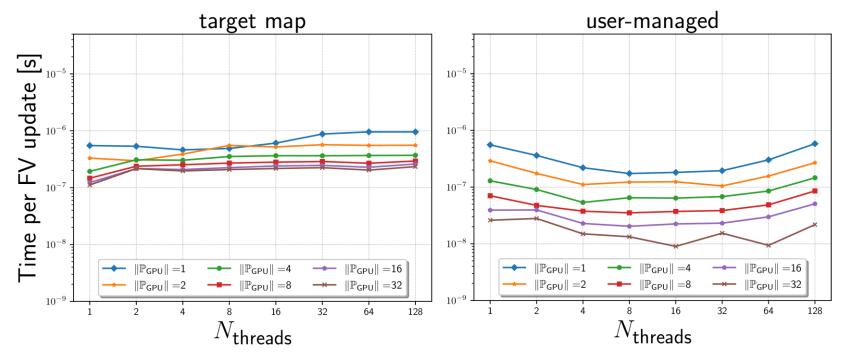


OpenMP Runtime Overhead with target map

- We do not constrain which core can access a GPU → multiple cores may hit the GPU simultaneously
- OpenMP runtime **locks** access to the GPU to prevent race conditions
- GPU can only read from page-locked host-pinned memory → requires **staging**
- We identify significant overhead and synchronisation
- Challenges:
 - Overlap computation and data transfer
 - Avoid data allocations/frees during computations



Overcoming the OpenMP Runtime Overhead



Realisation with target map

```
Procedure offload_map(||P<sub>GPU</sub>||, host_patch_data): // Data per patch are stored in one large array of structures (AoS)
1
   mapped_pointers \leftarrow allocate_host(||P_{GPU}||)
2
3
   for i \leftarrow 0 to ||P_{GPII}|| do
      patch_data \leftarrow host_patch_data[i]
4
5
      #pragma omp target enter data map(to:patch_data) // Map each patch's data onto the device
      6
7
   end
  #pragma omp target teams distribute map(to:mapped_pointers)
8
  for i \leftarrow 0 to ||P_{GPU}|| do
9
       // Do computations on Finite Volumes
10
11 end
12 for i \leftarrow 0 to ||P_{GPU}|| do
13
    patch data ← host patch data[i]
      #pragma omp target exit data map(from:patch_data) // Copy back the GPU outcomes to the host
14
15 end
16 mapped_pointers \leftarrow free_host()
```

Our Proposal: A GPU Memory Manager

- Each rank holds one instance of a thread-safe GPU memory manager
- **Reserves** memory on the GPU on demand
- Hands out **pre-reserved** memory to threads and reuses the memory
- Memory ownership resides on host
- Synchronisation and coordination of memory can be handled on the host
- Avoid locking by the OpenMP runtime
- Allocation routine returns a device pointer
- Free routine releases the device memory
- Share pre-allocated data between threads

Realisation with our GPU Memory Manager

- 1 **Procedure** offload_managed($||P_{GPU}||$, host_patch_data):
- 2 patch_data \leftarrow GPUMemoryManager \rightarrow allocate_device($||P_{GPU}||$) // Allocate GPU data and store the memory address
- 3 patch_data \leftarrow omp_target_memcpy(host_patch_data, $||P_{GPU}||$)

```
4 #pragma omp target teams distribute is_device_ptr(patch_data)
5 for i \leftarrow 0 to ||P_{GPU}|| do
```

6 // Do computations on Finite Volumes

7 end

8 host_patch_data \leftarrow omp_target_memcpy(patch_data, $||P_{GPU}||$)

9 patch_data ← GPUMemoryManager→free() // Release memory handle for possible reuse



Loop Collapsing and Reordering

```
1
  Procedure collapse(offset, N):
  #pragma omp target teams distribute parallel for simd collapse(3)
2
5
  for i \leftarrow 0 to N do
    for j \leftarrow 0 to N do
6
7
      for k \leftarrow -offset to N+offset do
        8
9
       end
10
    end
11 end
```



Loop Collapsing and Reordering

```
1
   Procedure collapse(offset, N):
   #pragma omp target teams distribute parallel for simd collapse(3)
2
5
   for i \leftarrow 0 to N do
     for j \leftarrow 0 to N do
6
7
        for k \leftarrow 0 to N+2. offset do
           volume index \leftarrow get_volume_index(i, j, k-offset)
8
9
        end
10
      end
11 end
```



Two Flavours of Equations

Euler Equations	Conformal and Covariant Z4 (CCZ4) Equations
 Describes the evolution of density, energy and velocity System of N = d + 2 non-linear partial differential equations d = 2 or 3 Low arithmetic intensity 	 Describes the evolution of space-time curvature Models gravitational waves System of N = 59 equations d = 3 only High arithmetic intensity
Solution via Finite Volumes with a generic Rusanov Riemann solver	

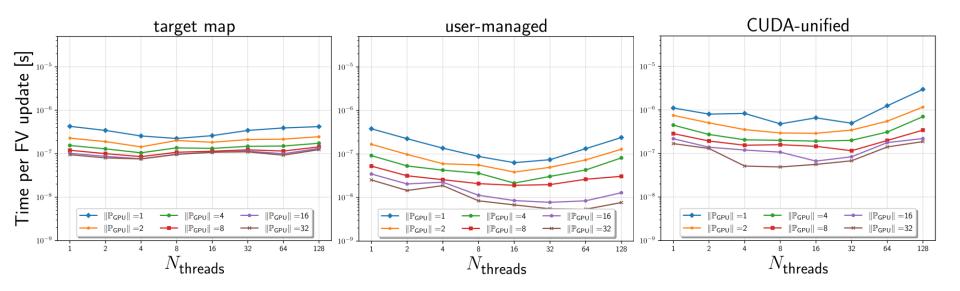


Benchmark Systems

Alex cluster: Erlangen National High Performance Computing Center (NHR @FAU)	JURECA DC Evaluation Platform: Jülich Supercomputing Centre (JSC)
 AlmaLinux 8.7 (Stone Smilodon) 2 × AMD EPYC 7713 Milan CPUs (64 cores per chip) 8 × NVIDIA A100 GPUs (80 GB memory) NVIDIA HPC SDK (v23.1) CUDA (v12.0) 	 Rocky Linux 8.7 (GreenObsidian) 2 × AMD EPYC 7443 Milan CPUs (24 cores per chip and SMT-2) 4 × AMD Instinct MI250 GPUs (128 GB memory, MCM) ROCm AOMP (v17.0-0)

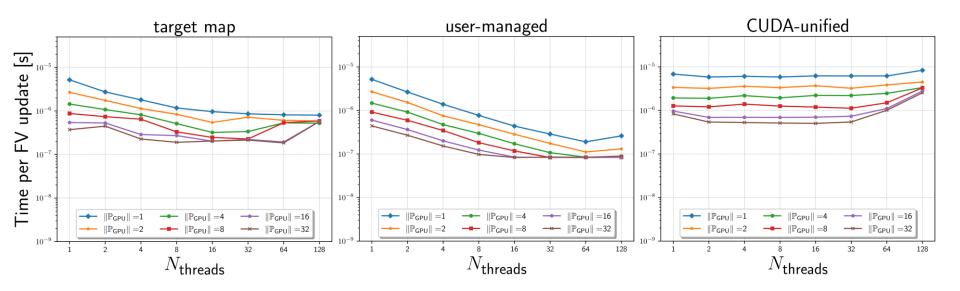


Performance Results – Euler Equations (NVIDIA)



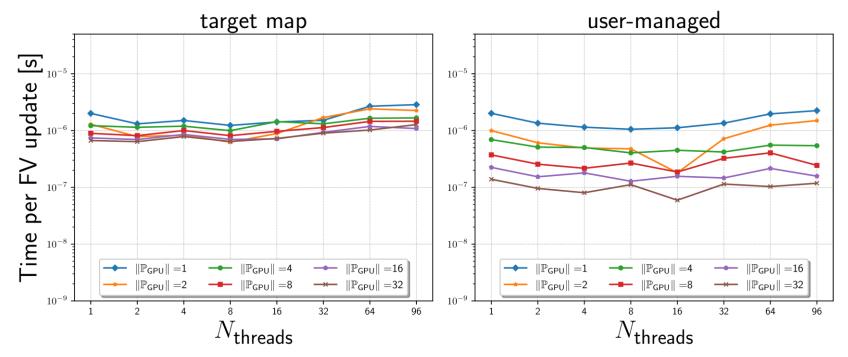


Performance Results – CCZ4 Equations (NVIDIA)



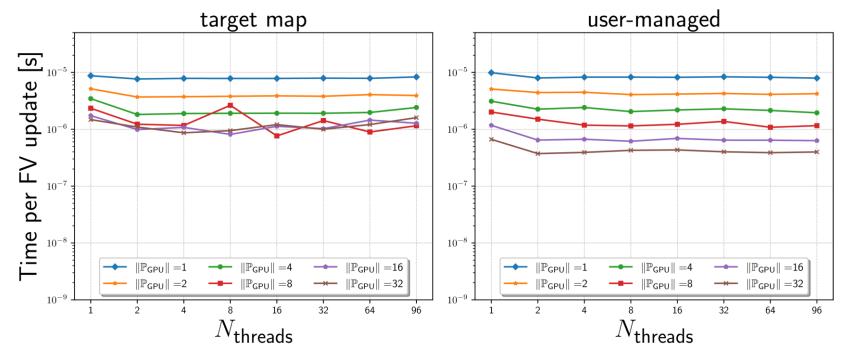


Performance Results – Euler Equations (AMD)





Performance Results – CCZ4 Equations (AMD)





Conclusion

- GPU Memory Manager withdraws memory management from the accelerator and assigns it to the host
- Host-centric realisation leads to a speedup of some calculations by an order of magnitude
- Increases robustness of concurrent offloading of patches
- No need for a static offloading pattern of patches or huge patches
- Presented a worst-case scenario motivated by a real-world science case

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Contact

- Feel free to contact me via Mail: <u>mario.wille@tum.de</u>
- Event platform Swapcard: Mario Wille
- Source code webpage: <u>www.peano-framework.org</u>
- Source code repository: <u>https://gitlab.lrz.de/hpcsoftware/Peano</u>
- Artifact description: <u>https://doi.org/10.5281/zenodo.7741217</u>

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Abstract. We identify and show how to overcome an OpenMP bottleneck in the administration of GPU memory. It arises for a wave equation solver on dynamically adaptive block-structured Cartesian meshes, which keeps all CPU threads busy and allows all of them to offload sets of patches to the GPU. Our studies show that multithreaded, concurrent, non-deterministic access to the GPU leads to performance breakdowns. since the GPU memory bookkeeping as offered through OpenMP's map clause, i.e., the allocation and freeing, becomes another runtime challenge besides expensive data transfer and actual computation. We, therefore, propose to retain the memory management responsibility on the host: A caching mechanism acquires memory on the accelerator for all CPU threads, keeps hold of this memory and hands it out to the offloading threads upon demand. We show that this user-managed, CPU-based memory administration helps us to overcome the GPU memory bookkeeping bottleneck and speeds up the time-to-solution of Finite Volume kernels by more than an order of magnitude.

Keywords: GPU offloading \cdot Multithreading \cdot OpenMP \cdot Dynamically adaptive mesh refinement

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